

The Newton polytope of implicit curves

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Outline

- 03. Toric elimination theory
- 15. Resultants
- 19. Implicitization
- 26. Implicit polygon

Toric elimination theory

Newton polytopes

The **support** A_i of a polynomial $f_i \in K[x_1^{\pm 1}, \dots, x_n^{\pm 1}]$, s.t.

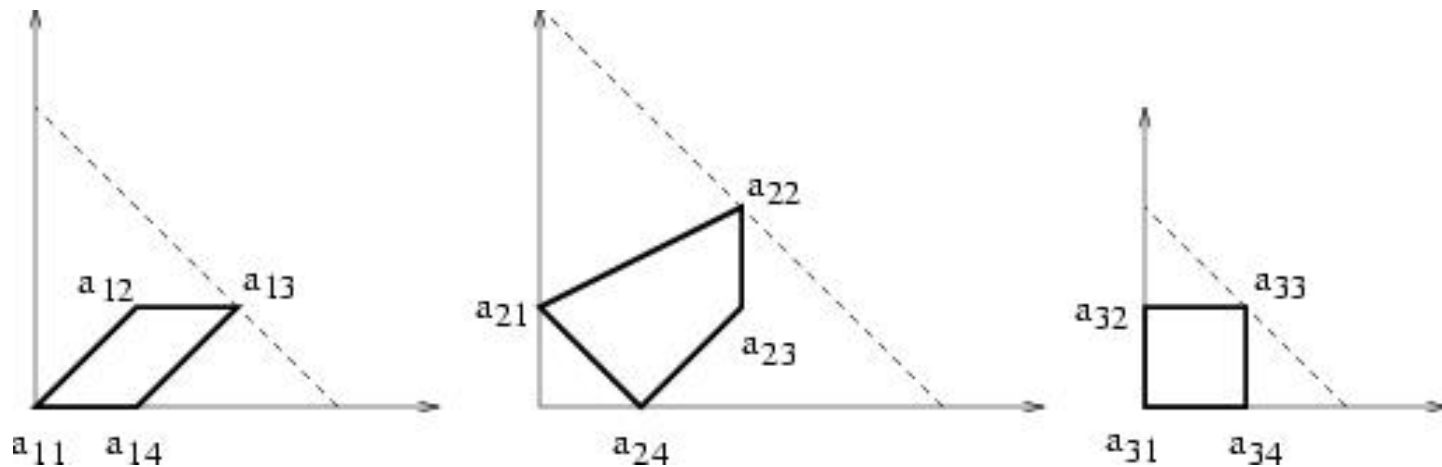
$$f_i = \sum_j c_{ij} x^{a_{ij}}, \quad c_{ij} \neq 0,$$

is defined as the set $A_i := \{a_{ij} \in \mathbb{Z}^n : c_{ij} \neq 0\}$.

The **Newton polytope** $Q_i \subset \mathbb{R}^n$ of f_i is the **Convex Hull** of all $a_{ij} \in A_i$.

Example:

$$\begin{aligned} f_1 &= c_{11} + c_{12}xy + c_{13}x^2y + c_{14}x \\ f_2 &= c_{21}y + c_{22}x^2y^2 + c_{23}x^2y + c_{24}x + c_{25}xy \\ f_3 &= c_{31} + c_{32}y + c_{33}xy + c_{34}x \end{aligned}$$



Mixed volume

1. The **mixed volume** $MV(P_1, \dots, P_n) \in \mathbb{R}$ of **convex** polytopes $P_i \subset \mathbb{R}^n$

- is **multilinear** wrt Minkowski addition and scalar multiplication:

$$MV(P_1, \dots, \lambda P_i + \mu P'_i, \dots, P_n) =$$

$$= \lambda MV(P_1, \dots, P_i, \dots, P_n) + \mu MV(P_1, \dots, P'_i, \dots, P_n), \quad \lambda, \mu \in \mathbb{R},$$

- st. $MV(P_1, \dots, P_1) = n! \operatorname{vol}(P_1)$.

2. Equivalently, $\operatorname{vol}(\lambda_1 P_1 + \dots + \lambda_n P_n)$ is a **polynomial** in scalar variables $\lambda_1, \dots, \lambda_n$, with **multilinear term** $MV(P_1, \dots, P_n) \lambda_1 \cdots \lambda_n$.

Mixed subdivisions

Regular subdivisions

For $Q_i \subset \mathbb{R}^n$, $(Q_i)_{i \in I} \rightarrow Q = \sum_{i \in I} Q_i : (q_i)_{i \in I} \mapsto \sum_{i \in I} q_i$.

Consider **lifting** functions $l_i : \mathbb{R}^n \rightarrow \mathbb{R}$, which define

$$\hat{Q}_i := \text{CH}\{(p_i, l_i(p_i)) : q_i \in Q_i\} \subset \mathbb{R}^{n+1}.$$

Let \hat{Q} be the Minkowski sum $\sum_i \hat{Q}_i$.

\forall face in the **lower-hull** of \hat{Q} is written uniquely as $\sum_i \hat{F}_i$, for faces $\hat{F}_i \subset \hat{Q}_i$.

\hat{Q} projects onto Q , so the lower-hull faces induce a **regular** subdivision of Q , with faces (cells) $\sum_i F_i$.

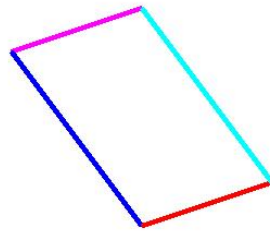
In particular, facets on the lower-hull project to maximal cells (dim = n).

Coherent subdivisions

Induced subdivisions are coherent,

i.e. there is a continuous change of the unique expression of every cell as we move to its subcells and adjacent cells.

We also say that the cells **intersect properly** as Minkowski sums.



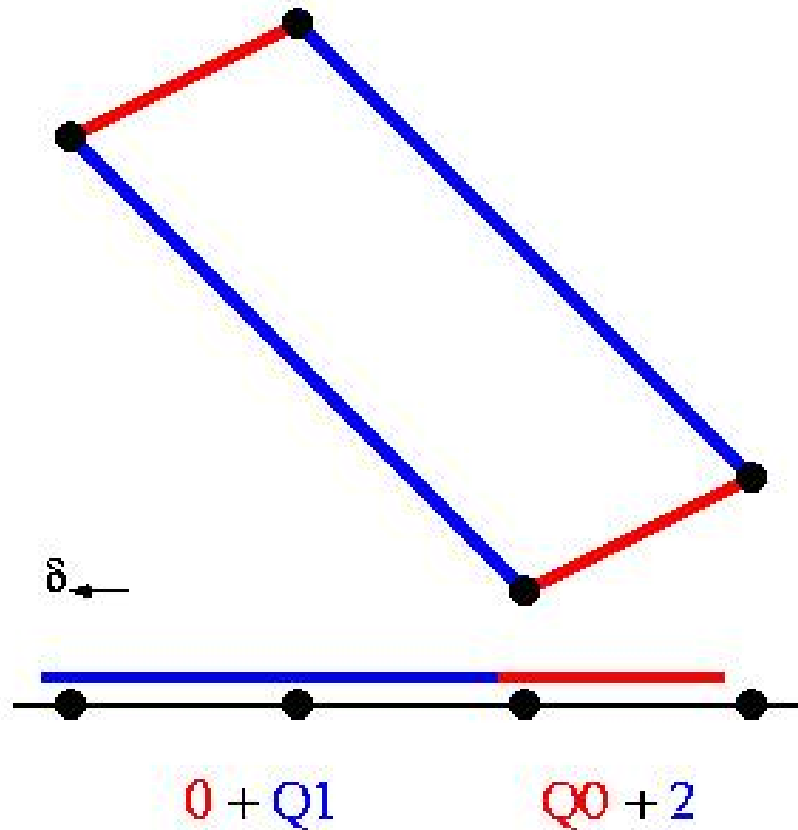
Eg: **In**coherent subdivision  of $Q_0 + Q_1$, $Q_i = [0, 1]$.

Leftmost cell = $\text{proj}(\hat{0} + \hat{Q}_1)$, so $\hat{0} + \hat{1} \mapsto 1 \in \mathbb{R}$.

Rightmost cell = $\text{proj}(\hat{1} + \hat{Q}_1)$, so $\hat{1} + \hat{0} \mapsto 1$: different expression.

Lifting in the Sylvester case

$$f_0 = c_{00} + c_{01}x, \quad f_1 = c_{10} + c_{11}x + c_{12}x^2$$



Point $2 = 0 + 2$ from both maximal cells.

Tight coherent mixed subdivisions

In general: $\dim(\sum_i F_i) \leq \sum_i \dim F_i$.

A **generic** lifting implies equality, i.e. a **tight** subdivision.

In particular, for a maximal-dimension cell, $n = \sum_i \dim F_i$.

Also, the lower-hull of \hat{Q} corresponds bijectively to Q .

E.g: **NOT** tight subdivision: 2 segments lifted in parallel:

$$\dim(F_0 + F_1) = 1 < \dim F_0 + \dim F_1 = 1 + 1.$$

One computes a (tight coherent) **mixed subdivision**, which partitions Q .

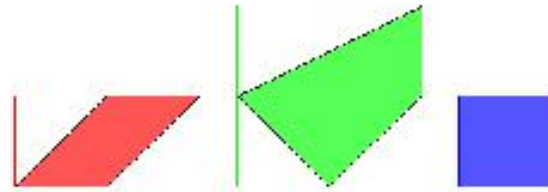


Figure 1: The given polytopes.

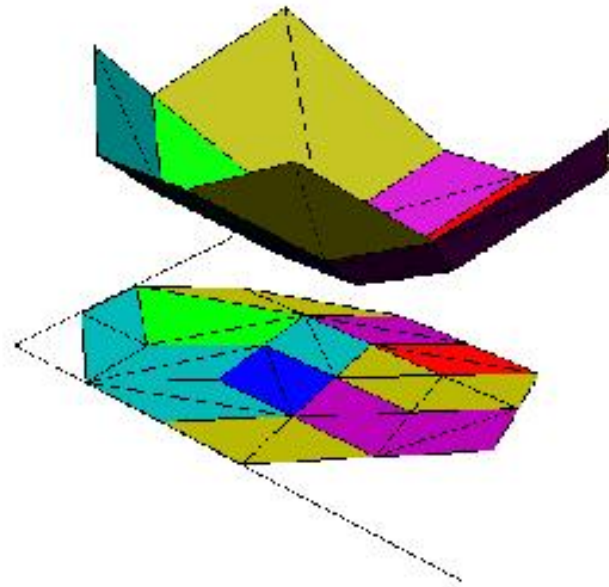


Figure 2: The lower hull of the lifted Minkowski Sum and its planar projection.

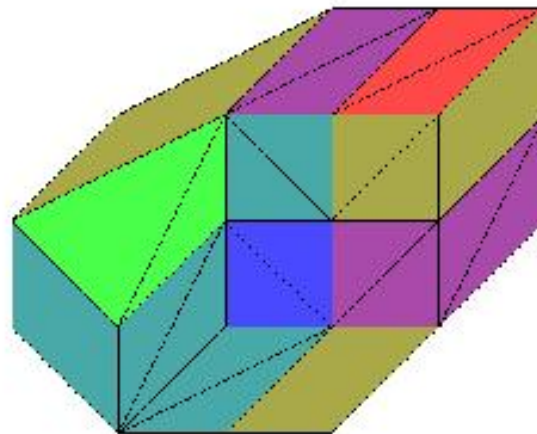


Figure 3: The mixed subdivision.

Example
lifting for
the over-
constrained
problem

Mixed cells

A maximal cell σ , in a mixed decomposition Δ , is **mixed** iff it has n linear summands, ie. n edges F_i : $\dim F_i = 1$.

- n polytopes: $Q = Q_1 + \dots + Q_n$, mixed cells are sums of edges.

Thm: $MV(Q_1, \dots, Q_n) = \sum_{\sigma} \text{vol}(\sigma)$, over all **mixed cells** $\sigma \in \Delta$.

- $n + 1$ polytopes: $Q = Q_0 + Q_1 + \dots + Q_n$, i -mixed cells are sums of edges plus vertex $a_i \in Q_i$.

Thm: $MV(Q_0, \dots, Q_{i-1}, Q_{i+1}, \dots, Q_n) = \sum_{\sigma} \text{vol}(\sigma)$,

over all **i -mixed cells** $\sigma \in \Delta$.

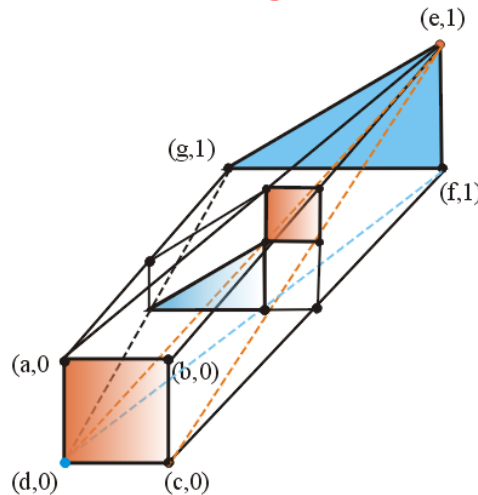
Enumerating mixed subdivisions

The **Cayley trick** introduces point set $C \subset \mathbb{Z}^{2n}$:

$$C := \{(e_i, a_{ij}) : i = 0, \dots, n, a_{ij} \in A_i\},$$

where the e_i constitute an affine basis of \mathbb{N}^n . So $|C| = |A_0| + \dots + |A_n|$.

Theorem. The set of all mixed subdivisions of $A_0, \dots, A_n \subset \mathbb{Z}^n$ corresponds bijectively to the set of all **regular triangulations** of C .



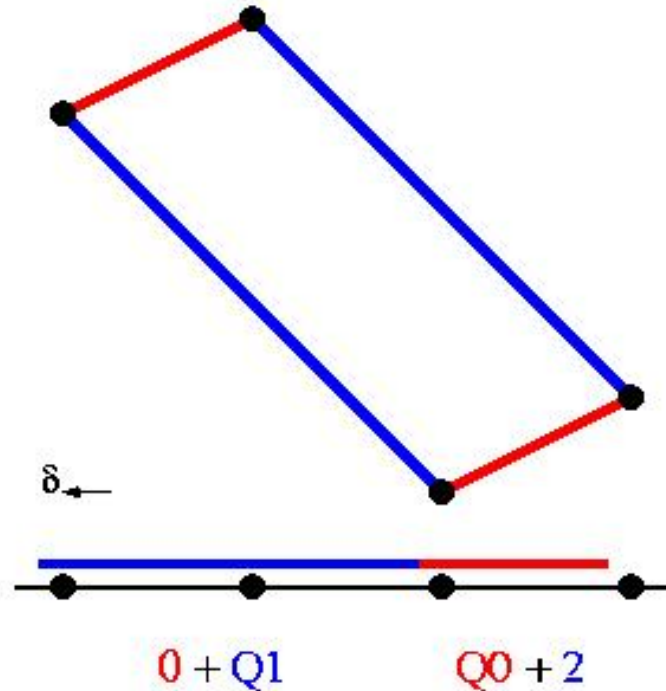
Caley trick in sparse Sylvester case

$$f_0 = c_{00} + c_{01}x,$$

$$f_1 = c_{10} + c_{12}x^2.$$

Cayley point set

$$C := \begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 2 \end{bmatrix}.$$



Triangulations: $\left(\begin{bmatrix} 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 2 \end{bmatrix} \right), \left(\begin{bmatrix} 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 2 \end{bmatrix} \right)$ shown,

and also $\left(\begin{bmatrix} 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \end{bmatrix} \right), \left(\begin{bmatrix} 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 2 \end{bmatrix} \right).$

Resultants

Resultant definition

Given $n + 1$ **Laurent** polynomials $f_0, \dots, f_n \in K[x_1, \dots, x_n, x_1^{-1}, \dots, x_n^{-1}]$ with indeterminate coefficients \vec{c} , their **projective**, resp. **toric / sparse**, *resultant* is the unique (up to sign) irreducible polynomial $R(\vec{c}) \in \mathbb{Z}[\vec{c}]$ such that

$$R(\vec{c}) = 0 \Leftrightarrow \exists \xi = (\xi_1, \dots, \xi_n) \in X : f_0(\xi) = \dots = f_n(\xi) = 0$$

where the variety X equals:

- the projective space \mathbb{P}^n over the algebraic closure \overline{K} ,
- resp. the **toric variety** X , $(\overline{K}^*)^n \subset X \subset \mathbb{P}^N$.

[van der Waerden, Gelfand-Kapranov-Zelevinsky, Cox-Little-O'Shea]

Newton polytope of the toric resultant

Given are supports A_0, \dots, A_n s.t. $k := \sum_i |A_i|$ and $\dim(\sum_i A_i) = n$. Consider the toric resultant $R \in \mathbb{Z}[c]$ and its Newton polytope in \mathbb{R}^k .

Let lifting $l \in \mathbb{R}^k$ define a (tight coherent) mixed subdivision of $Q_0 + \dots + Q_n$. Consider the **trailing monomial** of R with respect to l , which corresponds to the vertex of $\text{supp}(R) \subset \mathbb{Z}^k$ with inner normal l .

Theorem. [Sturmfels'94] This trailing monomial is

$$\prod_{i=0}^n \prod_{i\text{-mixed } \sigma} \text{coef}(f_i, a_i)^{\text{vol}(\sigma)},$$

where $\text{vol}(\cdot)$ denotes Euclidean volume and the **i -mixed cells** are $\sigma = F_0 + \dots + a_i + \dots + F_n : \dim a_i = 0$.

Newton polytope of the toric resultant (cont'd)

Corollary. A surjection exists from the set of **mixed-cell configurations** onto the **extreme monomials** of R (vertices of its Newton polytope).

Corollary. The coefficient of the trailing term is ± 1 .

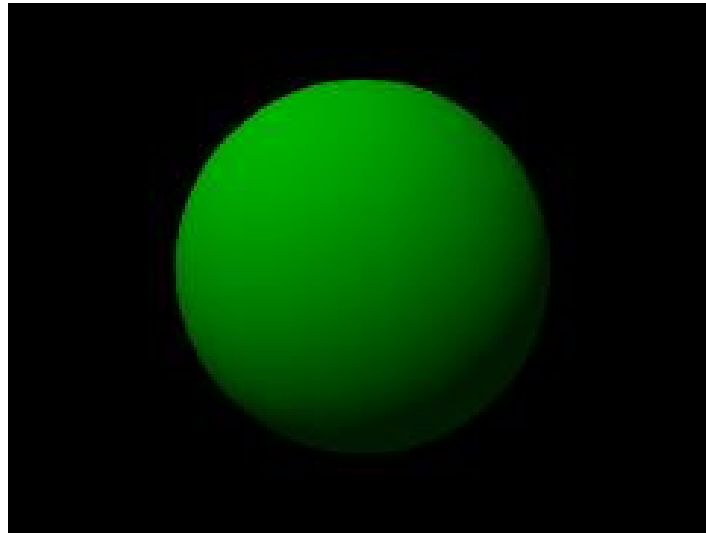
Implicitization

Example: sphere

The sphere in \mathbb{R}^3 is the set of **values** (x, y, z) :

$$x = \frac{t_1^2 - t_2^2 - 1}{t_1^2 + t_2^2 + 1}, y = \frac{2t_1}{t_1^2 + t_2^2 + 1}, z = \frac{2t_1 t_2}{t_1^2 + t_2^2 + 1}, t_1, t_2 \in [0, 1],$$

as well as the set of **roots** of $H(x, y, z) := x^2 + y^2 + z^2 - 1 = 0$.



Modeling/CAD use **parametric** and **implicit/algebraic** representations
 \Rightarrow must implicitize a (hyper)surface given a (rational) parameterization

Implicitization by linear algebra

S = monomials forming (a superset of) the **implicit support**.

C = unknown **coefficients** of implicit equation wrt S , $|C| = |S|$.

- $MC = \vec{0}$, where matrix M is $|S| \times |S|$, and contains values of S at points $(s_i, t_i), i = 1, \dots, |S|$. Try roots of unity.

- $(SS^T)C = \vec{0}$, substitute x, y, z by parametric expressions in $K[s, t]$, integrate over s, t ; solve for C [Corless-Galligo-Kotsireas-Watt'00].

Example: $\text{supp}(H) \subset \{x^3y, x^3, x^3y^2, y^2z^3\}$, then

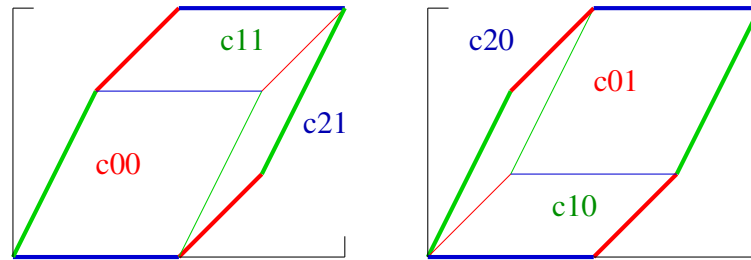
$$SS^T = \begin{bmatrix} x^6y^2 & x^6y & x^6y^3 & x^3y^3z^3 \\ x^6y & x^6 & x^6y^2 & x^3y^2z^3 \\ x^6y^3 & x^6y^2 & x^6y^4 & x^3y^4z^3 \\ x^3y^3z^3 & x^3y^2z^3 & x^3y^4z^3 & y^4z^6 \end{bmatrix} \Rightarrow C = \begin{bmatrix} -2 \\ 1 \\ 1 \\ -1 \end{bmatrix}.$$

- Approximate implicitization [Dokken].

Enumerate mixed subdivisions

A sparse example [Buchberger'88]

$$f_0 = c_{00} - c_{01}st, \quad f_1 = c_{10} - c_{11}st^2, \quad f_2 = c_{20} - c_{21}s^2.$$



The mixed subdivisions yield extreme monomials $c_{00}^4 c_{11}^2 c_{21}$, $c_{01}^4 c_{10}^2 c_{20}$.

The toric resultant turns out to be $R = c_{00}^4 c_{11}^2 c_{21} - c_{01}^4 c_{10}^2 c_{20}$.

The Fröberg-Dickenstein example

$$x = t^{48} - t^{56} - t^{60} - t^{62} - t^{63}, \quad y = t^{32}.$$

$$Q'_0 + 0, a + Q_1, Q''_0 + 32$$

$$\pm y^a c_{0a}^{32} c_{1,32}^{63-a}$$

$$a = 48, 56, 60, 62, 63$$

yields $\pm y^a$

$$a = 63$$

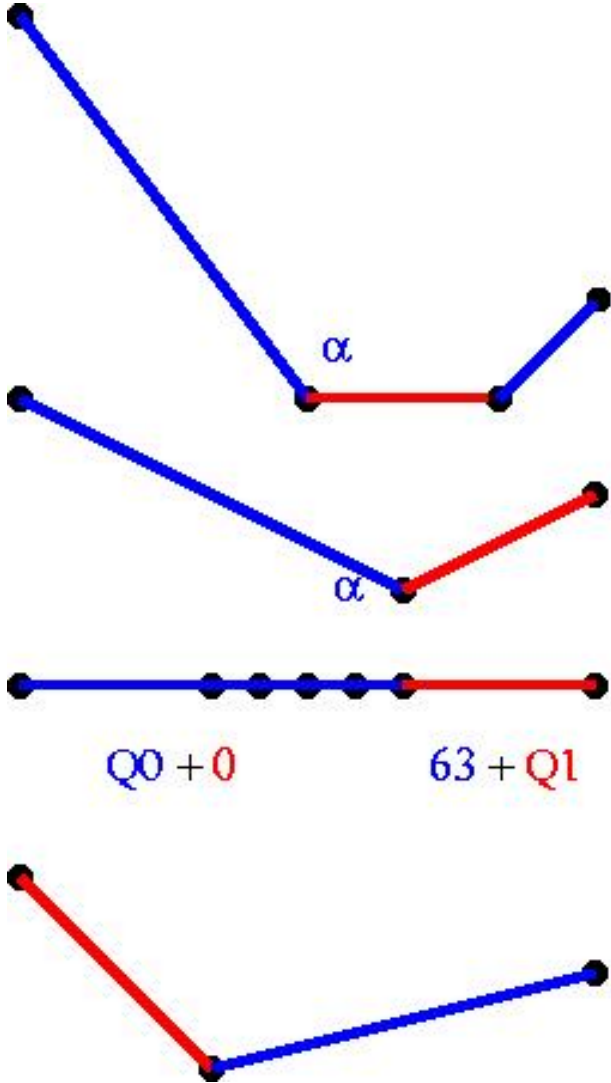
$$\pm y^{63} c_{0,63}^{32}$$

$$Q_0 + 0$$

$$63 + Q_1$$

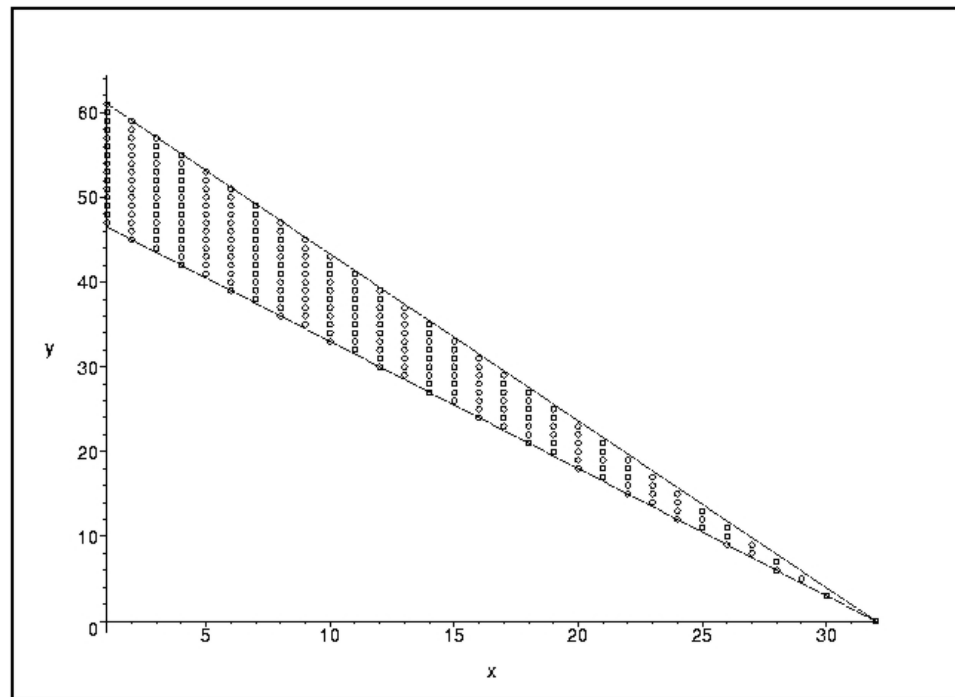
$$0 + Q_1, Q_0 + 32, (a = 0)$$

$$\pm x^{32} c_{1,32}^{63}$$



The Fröberg-Dickenstein example (cont'd)

The projected support is defined by points $(0, 48)$, $(0, 63)$, $(32, 0)$:
This triangle includes 257 lattice points, optimally:



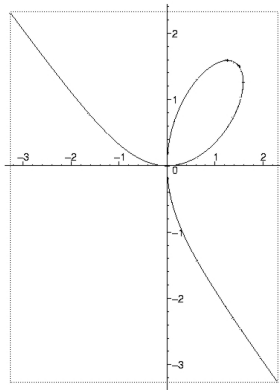
Degree bounds yield quad with additional vertices $(0, 0)$, $(32, 31)$.

Implicitization examples [E-Kotsireas'03]

[Descartes' folium]

[1596-1650]

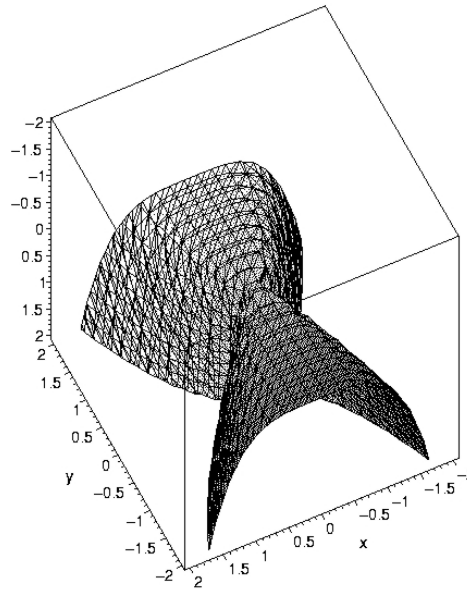
$$(x, y) = \left(\frac{3 t^2}{t^3 + 1}, \frac{3 t}{t^3 + 1} \right)$$



$$H = x^3 + y^3 - 3 x y$$

[Buchberger'88]

$$(x, y, z) = (st, st^2, s^2)$$



$$H = x^4 - y^2 z$$

[Busé'01]

$$x = \frac{s^2}{s^3 + t^3},$$

$$y = \frac{s^3}{s^3 + t^3},$$

$$z = \frac{t^2}{s^3 + t^3}$$

$$H = x^3 - 2x^3y + x^3y^2 - y^2z^3$$

Implicit polytope

Consider parameterizations with **fixed supports** and **generic** coefficients.

- Compute the resultant's Newton polytope, then specialize.
 - [E-Kotsireas'03] developed Maple code based on Topcom [Rambau].
 - [EKP'07] algorithm for projecting polytope to $\mathbb{R}, \mathbb{R}^2, \mathbb{R}^3$.
- Tropical geometry for the polytope of **Laurent-polynomial** hypersurfaces and varieties of $\text{codim} > 1$. For curves, determines implicit vertices [Sturmfels-Tevelev-Yu'06].
- Implicit edges for all curves, using **linear factors**:
 - [D'Andrea-Sombra'07] use mixed fiber polytopes [Esterov-Khovanskii'07], characterize polygons which realize as implicit.
 - [Dickenstein-Sturmfels'07] use tropical discriminants.
- Specify implicit vertices for all curves [E-Konaxis-Palios'07]

Polynomially parameterized curves

$$x_i = P_i \in K[t], \quad i = 0, 1,$$

with supports $\{a_0, \dots, a_n\}, \{b_0, \dots, b_m\} \subset \mathbb{N}$.

Cayley's trick: $\{(a_0, 0), \dots, (a_n, 0), (b_0, 1), \dots, (b_m, 1)\} \subset \mathbb{N}^2$.

Now, consider **triangulations** of this point set.

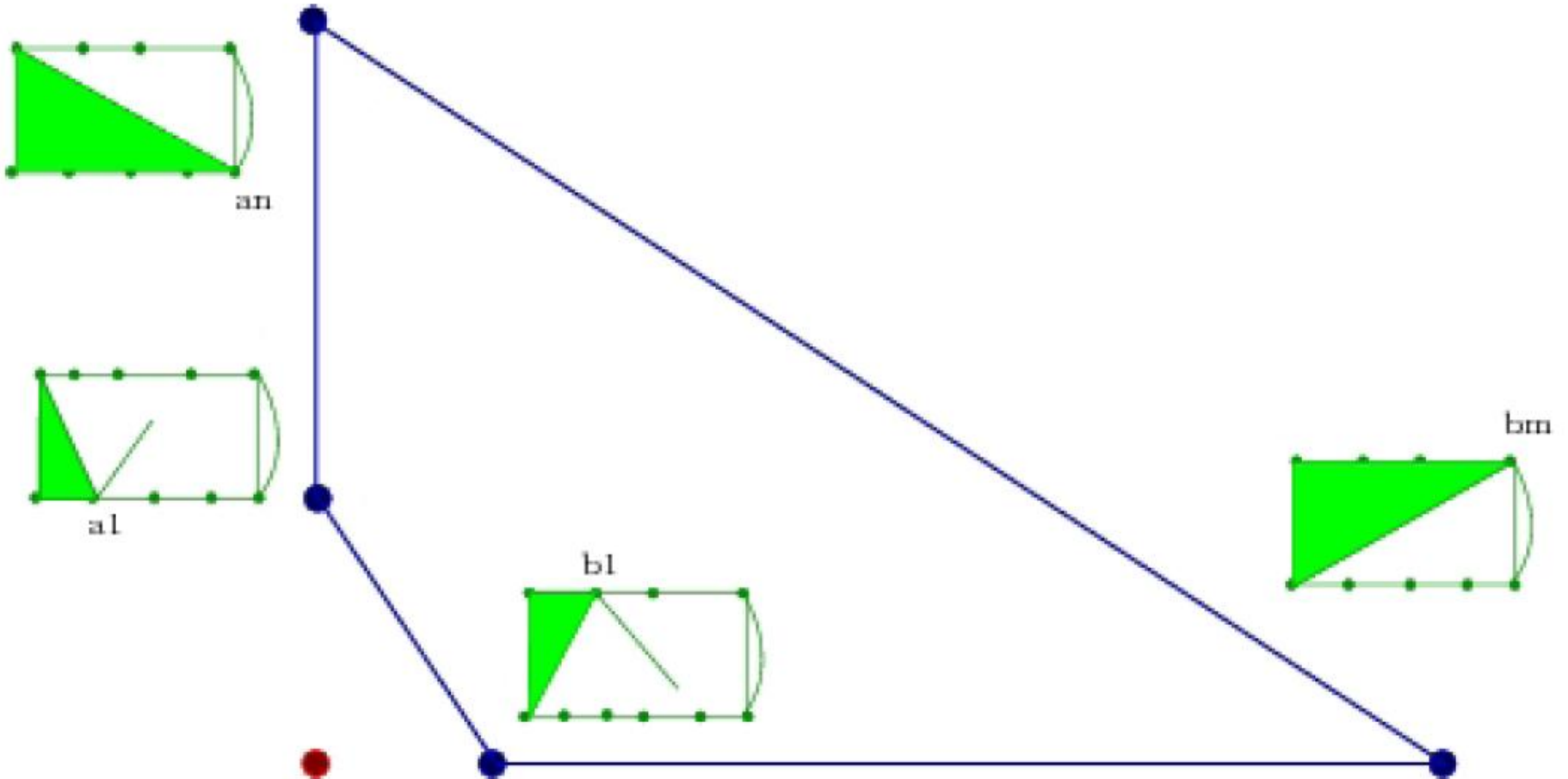
Triangle $(a_0, b_i, b_j) \rightarrow x^{b_j - b_i}, (a_i, a_j, b_0) \rightarrow y^{a_j - a_i}$.

Triangles $(a_i, *, b_j), i, j \neq 0$, do not contribute to x, y .

We shall say a_0, b_0 are **selected**.

Implicit polygon of Polynomial curves

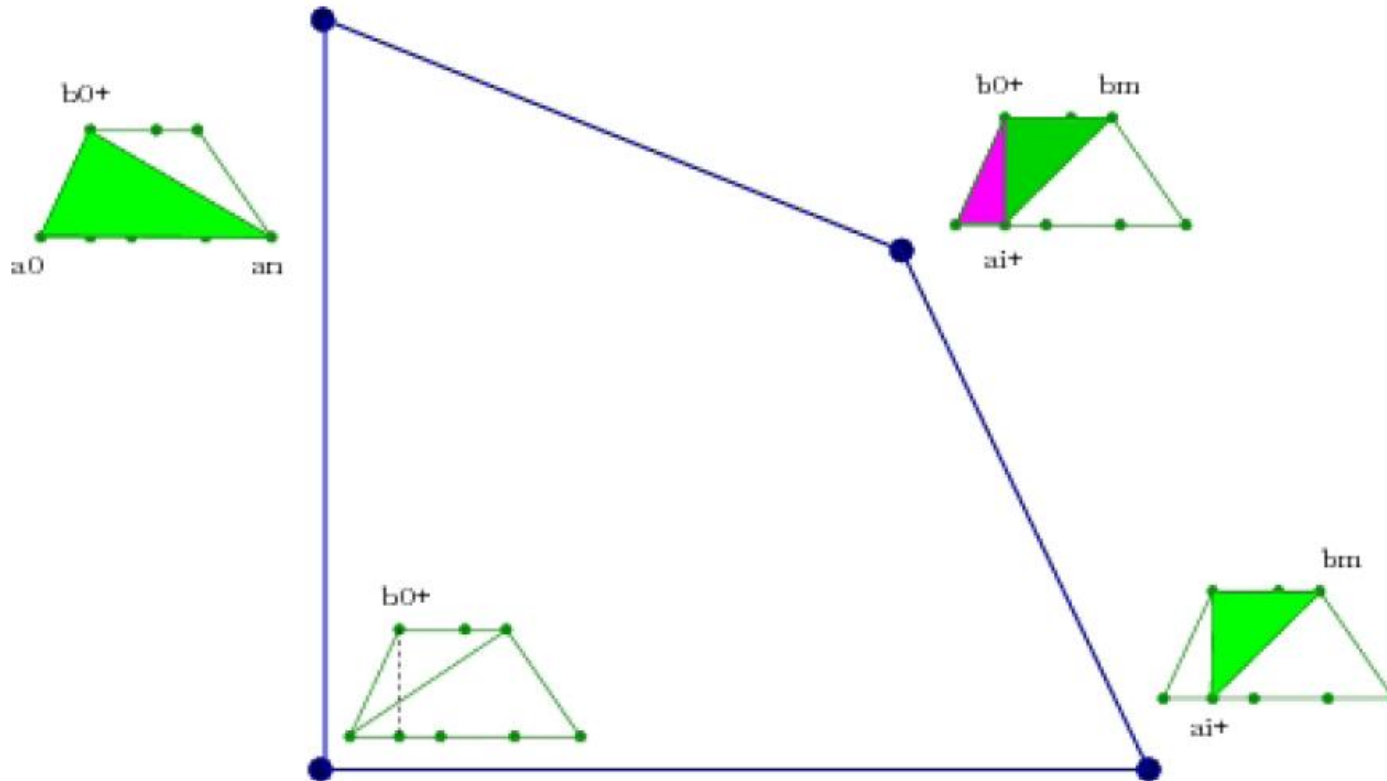
\exists constant term in some $P_i(t) \Rightarrow \exists$ implicit vertex $(0, 0)$:



Cor: $\text{coef}(x^{b_m}) = \pm(-c_{1_m})^{a_n}$, $\text{coef}(y^{a_n}) = \pm(-c_{0_n})^{b_m}$.

Laurent-polynomial parameterization

$\{a_0, \dots, a_n\}, \{b_0, \dots, b_m\} \subset \mathbb{Z}$, unique selected a_i^+, b_0^+ .



Up-right vertex = $(b_m, |a_0|)$ iff $\det \begin{bmatrix} |a_0| & a_n \\ |b_0| & b_m \end{bmatrix} > 0$, $(|b_0|, a_n)$ iff $\det < 0$.

Rational curves, different denominators

$$x_i = \frac{P_i(t)}{Q_i(t)}, \gcd(P_i, Q_i) = 1 \rightarrow f_i = x_i Q_i(t) - P_i(t) \in K[t], \quad i = 0, 1,$$

where $\text{supp}(f_0) = \{a_0, \dots, a_n\}$, $\text{supp}(f_1) = \{b_0, \dots, b_m\} \subset \mathbb{N}$.

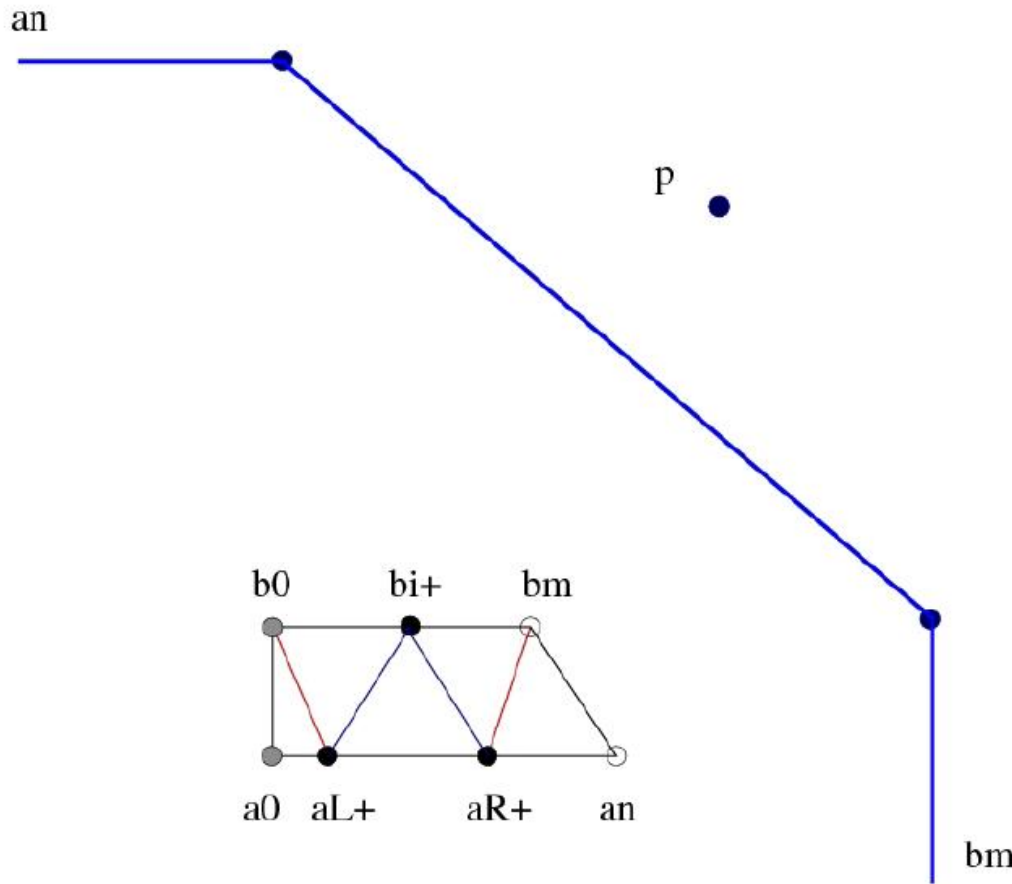
Lemma. Consider direction $(1, 1)$. The **upper** [resp. **lower**] hull of the Newton polygon has vertices of the form:

$$\left(\sum_{k,l,r} \text{vol}(a_k^+, b_l, b_r), \sum_{l,r,j} \text{vol}(a_l, a_r, b_j^+) \right) \in \mathbb{N}^2,$$

where $a_k^+ \in \text{supp}(Q_0)$ [resp. $\text{supp}(Q_0) - \text{supp}(P_0)$],
and $b_j^+ \in \text{supp}(Q_1)$ [resp. $\text{supp}(Q_1) - \text{supp}(P_1)$]

Proof. The significant resultant coefficients are the following:
the monomials $c_i x_i$ and the binomials $c_i x_i + c'_i \in K[x_i]$
[resp. the monomials $c_i x_i \in K[x_i]$].

Upper-right corner



$$x_{\max} = b_m,$$

$$y_{\max} = (a_R^+ - a_L^+) + \mathcal{X}(b_m^+) \cdot (a_n - a_R^+)$$

if $a_0^-, b_0^-, a_n^-, b_m^-$ then:

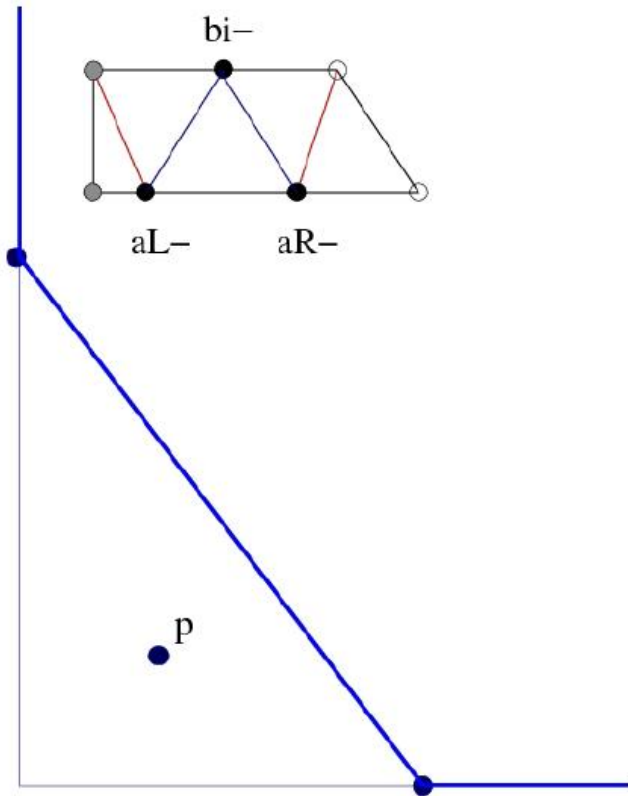
$$\exists p = (b_R^+, a_n - a_L^+) \Leftrightarrow$$

$$\det \begin{bmatrix} a_n - a_R^+ & a_L^+ \\ b_L^+ & b_m - b_R^+ \end{bmatrix} > 0,$$

$$p = (b_m - b_L^+, a_R^+) \Leftrightarrow \det < 0$$

Selected $a_k^+ \in \text{supp}(Q_0), b_j^+ \in \text{supp}(Q_1)$, not selected a_k^-, b_j^- .

Lower-left corner



$$x_{\min} = 0,$$

$$y_{\min} = \mathcal{X}(b_0^+)a_L^+ + \mathcal{X}(b_m^+)(a_n - a_R^-)$$

if $a_0^+, b_0^+, a_n^+, b_m^+$ then:

$$\exists p = (b_L^-, a_n - a_R^-) \Leftrightarrow$$

$$\det \begin{bmatrix} a_n - a_R^- & a_L^- \\ b_m - b_R^- & b_L^- \end{bmatrix} < 0,$$

$$p = (b_m - b_R^-, a_L^-) \Leftrightarrow \det > 0$$

Selected $a_k^+ \in \text{supp}(Q_0) - \text{supp}(P_0)$, b_j^+ , not selected a_k^-, b_j^- .

Rational curves, common denominator

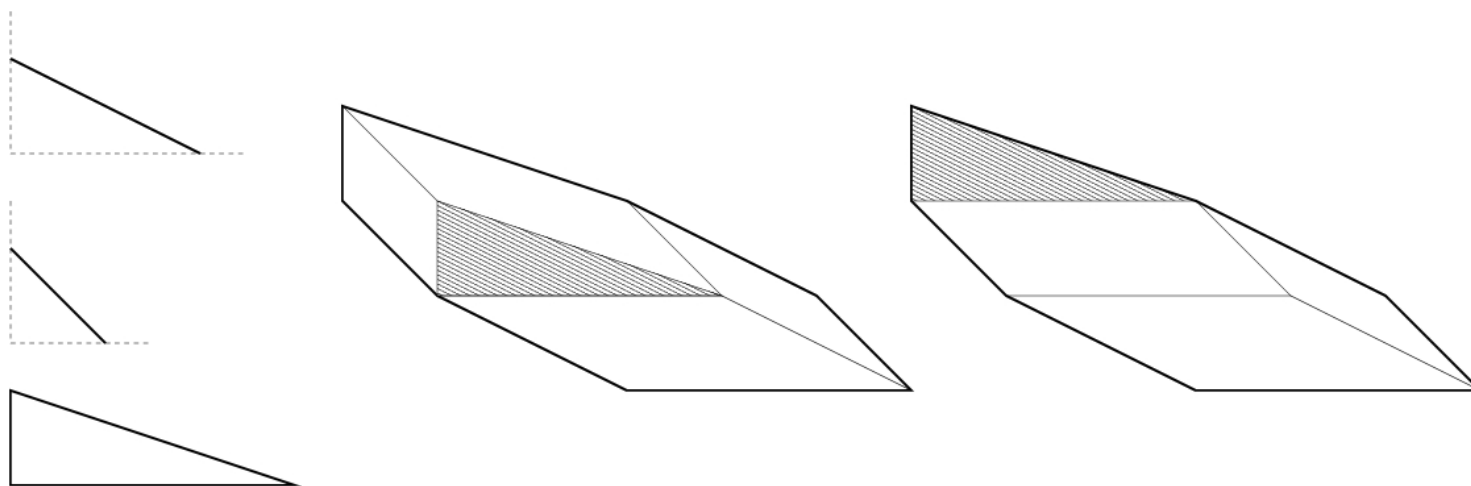
$$x_0 = \frac{P_0(t)}{Q(t)}, \quad x_1 = \frac{P_1(t)}{Q(t)}, \quad \gcd(P_i, Q) = 1,$$

$$B_i = \text{supp}(P_i) = \{b_{iL}, \dots, b_{iR}\}, \quad i = 0, 1, \quad B_2 = \text{supp}(Q) = \{b_{2L}, \dots, b_{2R}\} \subset \mathbb{N}$$

Let $f_0 = x_0 r - P_0(t)$, $f_1 = x_1 r - P_1(t)$, $f_2 = r - Q(t) \in K[t, r]$,

$$A_i = \text{supp}(f_i) = \{a_{i0} = (0, 1), a_{iL} = (b_{iL}, 0), \dots, a_{iR} = (b_{iR}, 0)\}, \quad i = 0, 1, 2$$

Example: Folium of Descartes: $x = \frac{3t^2}{1+t^3}$, $y = \frac{3t}{1+t^3}$.

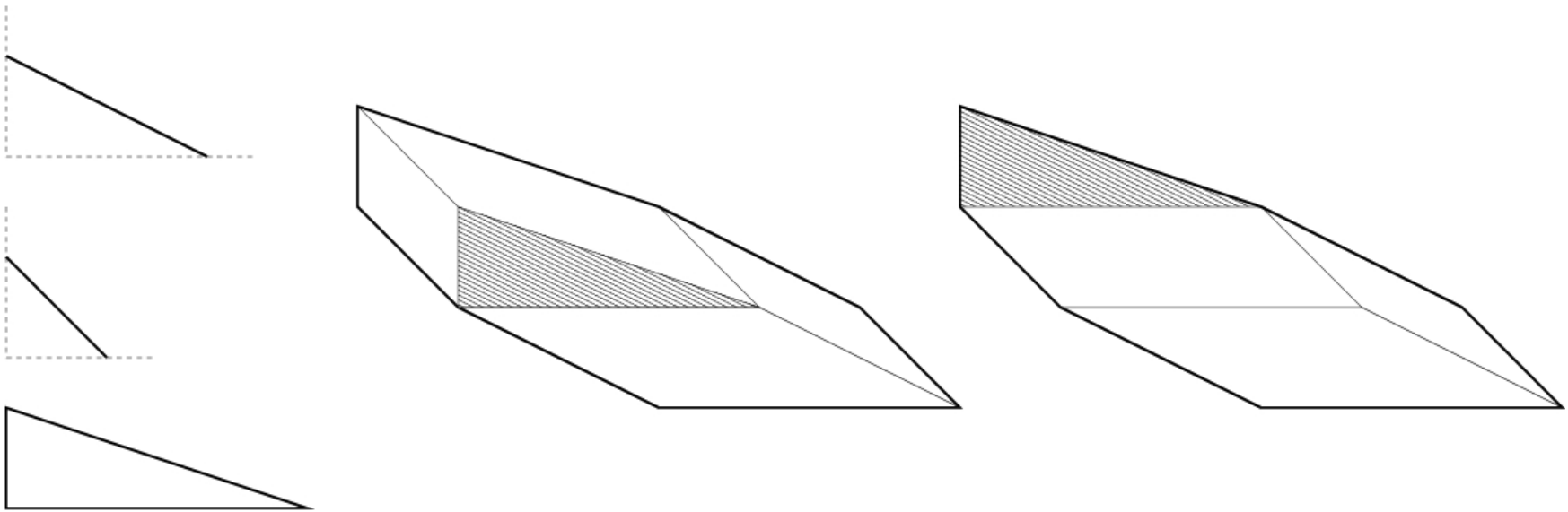


Mixed subdivisions

Lemma. In Minkowski sum $\text{CH}(A_0) + \text{CH}(A_1) + \text{CH}(A_2)$, all a_{i0} -mixed cells are as follows, for $\{i, j, k\} = \{0, 1, 2\}$:

$$a_{i0} + E_j + E_k, \text{ where edge } E_j = (a_{j0}, a_{jt}) \subset A_j,$$

and $E_k \subset A_k$ is either **non-horizontal** (a_{k0}, a_{km}) , or **horizontal** (a_{kl}, a_{km}) .



Cor. It suffices to consider subdivisions of **segment** $((0, 2), (u, 2))$.

Implicit equation

Lemma. Implicit equation $\phi \in K[x_0, x_1]$ has

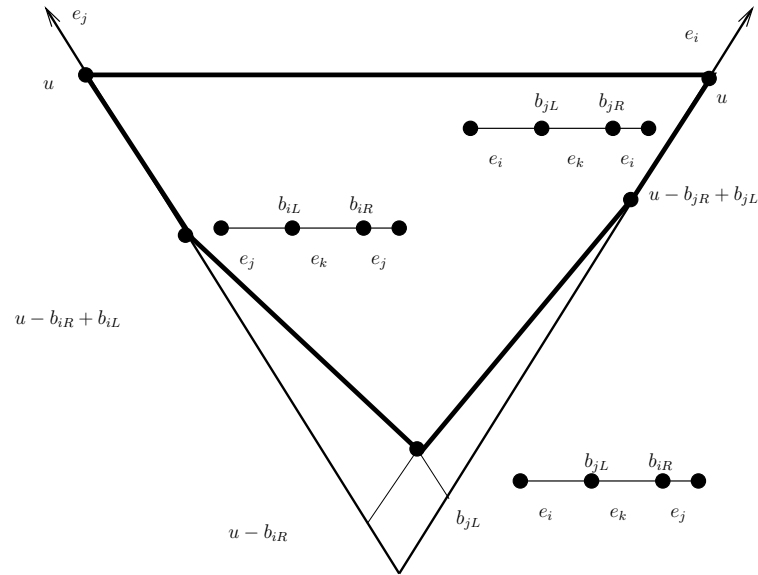
$$\deg \phi = \text{vol}[\text{CH}(\cup_i \text{supp}(f_i))] = \max_i \{b_{iR}\} - \min_i \{b_{iL}\} = u - 0.$$

Cor. $\text{Res}_{t,r}(f_0, f_1, f_2)$ is homogeneous, of degree u , in the 3 coefficients corresponding to the a_{00}, a_{10}, a_{20} .

$\Phi :=$ specialization of $\text{Res}_{t,r}(f_0, f_1, f_2)$ to polynomial in the 3 coeffs:
 $N(\Phi) \subset \mathbb{R}^3$ projects bijectively to $N(\phi) \subset \mathbb{R}^2$.

Thm. For computing the vertices of $N(\phi)$ (or $N(\Phi)$), it suffices to consider subdivisions defined by the **vertices** of the $\text{CH}(A_i), i = 0, 1, 2$, i.e. the relevant liftings of each A_i are linear in \mathbb{R}^2 .

Implicit vertices (A)



Thm. If all $\text{CH}(B_i \cup B_j) = [0, u]$, then $N(\phi) = \text{CH}((0, 0), (0, u), (u, 0))$.
 Otherwise, if $\exists! B_k = [0, u]$, then the (e_i, e_j) -vertices of $N(\phi)$ lie in

$$\{(u, 0), (0, u), (0, u - b_{iR} + b_{iL}), (b_{jL}, u - b_{iR}), (u - b_{jR} + b_{jL}, 0)\},$$

where $\{i, j, k\} = \{0, 1, 2\}$, and $b_{iL}(u - b_{jR}) \geq b_{jL}(u - b_{iR})$.

Implicit vertices (B)

Thm. If $\forall B_t \neq [0, u]$, then choose $\{i, j, k\} = \{0, 1, 2\}$ s.t.:

$$0 < b_{iL} \leq b_{iR} = u, 0 = b_{jL} \leq b_{jR} < u, 0 \leq b_{kL} \leq b_{kR} < u.$$

If $b_{kL} > 0$, the (e_i, e_j) -vertices lie in

$$\{(b_{jR}, 0), (b_{kR}, u - b_{kR}), (b_{kL}, u - b_{kL}), (0, u - b_{0L}), (0, 0)\}.$$

If $b_{kL} = 0$, the 3rd and 4th vertices are replaced by $(0, u)$.

Implicit polygon cuts

Corollary. Start with a triangle or quadrilateral that has a vertex at $(0, 0)$ and incident edges which lie on the axes.

- Polynomial parameterizations:

Take a right triangle, apply at most one corner cut excluding the origin.

- Rational parameterizations with equal denominators:

Take a right triangle, apply at most two cuts (same or different corners).

- Rational parameterizations with different denominators:

Take a quadrilateral, apply at most two cuts (same or different corners).

Conclusions

Inverse. If $\phi(x, y)$ admits a polynomial parameterization, then $N(\phi)$ has one edge on its upper hull wrt $(1, 1)$.

Then, if $N(\phi) = \text{segment}$, it contains no interior lattice points.

Future:

- Specify genericity conditions, determine extremal coefficients.
- Polytope of implicit Surfaces.
- When are linear liftings sufficient? Pyramids? Separated variables?
- Project resultant polytope to low dimension.
- Numerical implicitization.

Secondary polytopes

Secondary polytope

Consider the graph of **regular triangulations** of point-set $C \subset \mathbb{Z}^d$, where edges correspond to (bistellar) flips.

Theorem [Gelfand-Kapranov-Zelevinsky, Billera-Sturmfels]

If C affinely spans \mathbb{R}^d , then the graph can be embedded in $\mathbb{R}^{|C|-d-1}$ as the **secondary polytope** $\Sigma(C)$. For triangulation T ,

$$(v_T)_i = \sum_{i \in \text{vtx}(\sigma): \sigma \in T} \text{vol}(\sigma), \quad i = 1, \dots, |C|,$$

where $\text{vtx}(\sigma)$ are the vertices of simplex σ .

E.g. $C \subset \mathbb{Z}^2$, $|C| = 4$:



Circuits

A **circuit** $Z = \{c_1, \dots, c_t\}$ is a minimal affinely-dependent subset of C , satisfying $\lambda_1 c_1 + \dots + \lambda_t c_t = 0$, where $\lambda_i \neq 0$, $\sum_i \lambda_i = 0$.

Z admits triangulations $Z^+ = \{Z \setminus \{c_i\} \mid \lambda_i > 0\}$, $Z^- = \{Z \setminus \{c_i\} \mid \lambda_i < 0\}$. Each **flip** $T \leftrightarrow T'$ corresponds to precisely one circuit Z s.t.

$$T' \simeq T \setminus Z^+ \cup Z^-$$

E.g. $Z = C$,
$$-\begin{bmatrix} 0 \\ 0 \end{bmatrix} + \begin{bmatrix} 1 \\ 0 \end{bmatrix} + \begin{bmatrix} 0 \\ 1 \end{bmatrix} - \begin{bmatrix} 1 \\ 1 \end{bmatrix},$$



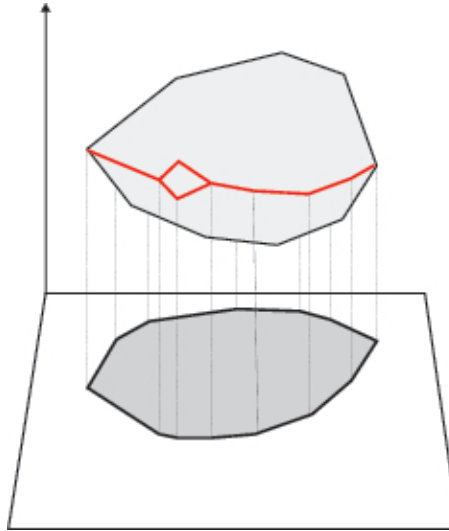
Projecting $\Sigma(C)$ to \mathbb{R}

Project to 1st coordinate, corresponding to $c_1 \in C \subset \mathbb{Z}^d$ (point set).

- Let T be a regular triangulation, and consider flip $T \leftrightarrow T'$. Then c_1 is a vertex of every new simplex iff $(v_T)_1 < (v_{T'})_1$.
- Let Z_j be the flips that make $(v)_1$ increase, and σ_j the unique simplex vanishing with Z_j not containing c_1 . Then, the triangulation T maximizing $(v_T)_1$ is s.t. the volume of simplices containing σ_j is max.
- Hence $(v)_1 \uparrow$; if strictly \uparrow then min-path.

[E-Konaxis-Palios'07]

Projecting $\Sigma(C)$ to $\mathbb{R}^k, k \geq 2$



- **Complexity:** Time = $O^*(s^2 m) \text{LP}(\dim \Sigma, s)$, Space = $O(ns)$,
 $s = \max \# \text{any-dim simplices} = O(k^n)$, $m = \# \text{mixed-cell config's}$.

- Gift-wrapping, $\text{CCW}(u, v, w) = \det \begin{bmatrix} 1 & u_1 & u_2 \\ 1 & v_1 - u_1 & v_2 - u_2 \\ 1 & w_1 - v_1 & w_2 - v_2 \end{bmatrix}$.

Goal: complexity proportional to $\# \text{silhouette-points}$.