

Variations on Diffie Hellman Key Exchange

Let the following key exchange protocol with input $\langle g, G, m \rangle$, which can be described by the following moves:

1. Alice selects $x_A \xleftarrow{R} \mathbb{Z}_m$ and sends $y_A = g^{x_A}$ to Bob.
2. Bob selects $r \xleftarrow{R} \mathbb{Z}_m$ and calculates $k_B = g^r$, while sending $z = y_A^r$ to Alice.
3. Alice calculates $k_A = z^{\frac{1}{x_A}}$, which is the key.

Correctness

$$k_A = z^{\frac{1}{x_A}} = ((g^{x_A})^r)^{\frac{1}{x_A}} = g^r = k_B$$

Security

We will show that the above described protocol is secure under the DDH assumption:

$$|\text{Prob}[B(\langle g, G, m \rangle, g^x, g^y, g^{xy}) = 1] - \text{Prob}[B(\langle g, G, m \rangle, g^x, g^y, g^z) = 1]| \leq \text{negl}(1^\lambda)$$

Suppose that

$$\text{Prob}_{key \leftarrow \text{Key}(\lambda)}[V(key) = 1] = \delta$$

meaning that the Adversary can extract some part of the key.

Suppose that the protocol is not secure. Then there exists an algorithm A and a predicate V such that

$$\text{Prob}[A(\tau) = V(key(\tau))] \geq \max\{\delta, 1 - \delta\} + \alpha \tag{1}$$

with α a non negligible function of λ . Then for $trans(1^\lambda) = \langle g, G, m, y, z \rangle$ we have that

$$key(\langle g, G, m, y, z \rangle) = z^{\frac{1}{\log_g y}}$$

Let

$$D_\lambda = \{\langle g, G, m \rangle \leftarrow GGen(1^\lambda); a, b \xleftarrow{R} \mathbb{Z}_m : (G, m, g^a, g^{ab}, g^b)\}$$

and

$$R_\lambda = \{\langle g, G, m \rangle \leftarrow GGen(1^\lambda); a, b, c \xleftarrow{R} \mathbb{Z}_m : (G, m, g^a, g^c, g^b)\}$$

We devise a statistical test B that on input $\langle g, G, m, a, b, c \rangle$ sends $(\langle g, G, m \rangle, a, c)$ to A , which extracts σ and if $V(b) = \sigma$, then B outputs 1, else B outputs 0.

Then from (1) we have that

$$\text{Prob}_{\gamma \leftarrow D_\lambda}[B(\gamma) = 1] \geq \max\{\delta, 1 - \delta\} + \alpha$$

and

$$\begin{aligned}\text{Prob}[B(\gamma) = 1]_{\gamma \leftarrow R_\lambda} &= \text{Prob}[A(\tau_\gamma) = 1] \cdot \text{Prob}[V(c) = 1] + \text{Prob}[A(\tau_\gamma) = 0] \cdot \text{Prob}[V(c) = 0] \\ &= \text{Prob}[A(\tau_\gamma) = 1]\delta + \text{Prob}[A(\tau_\gamma) = 0](1 - \delta) \\ &\leq \text{Prob}[A(\tau_\gamma) = 1] \max\{\delta, 1 - \delta\} + \text{Prob}[A(\tau_\gamma) = 0] \max\{\delta, 1 - \delta\} \\ &= \max\{\delta, 1 - \delta\}^*\end{aligned}$$

and

$$|\text{Prob}[B(D) = 1] - \text{Prob}[B(R) = 1]| = |\max\{\delta, 1 - \delta\} + \alpha - \max\{\delta, 1 - \delta\}| = \alpha$$

which we have supposed to be non negligible, and therefore leads us to a contradiction since we have assumed the DDH assumption.

*More accurately $\text{Prob}[V(c) = 1] = \delta'$ but it is easy to show that $|\delta - \delta'| = \text{negl}(\lambda)$