The Web Ontology Language OWL2

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Knowledge Technologies

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- As a DL-based language:
 - different language profiles based on expressiveness are possible
 - ontologies comprise a TBox, an RBox and ABox
 - sound and complete algorithms for the reasoning tasks are available

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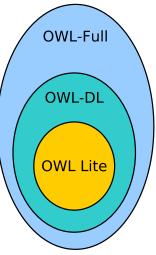
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- No Unique Name Assumption:
 - e.g., :Mary :hasChild :John ; :hasChild :kid1 :John and :kid1 may refer to the same entity
- Open World Assumption:
 - absence of information must not be considered as negative information, e.g. :Mary :hasChild :John does not entail that Mary has only one child

OWL Overview (cont'd)

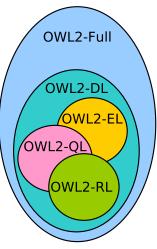
OWL profiles (sublanguages) are syntactic restrictions:

- **OWL Lite:** designed for easy implementation with a functional subset of OWL.
- **OWL DL:** designed to support the existing Description Logic and provide a language subset that has desirable computational properties for reasoning systems.
- **OWL Full:** relaxes some of the constraints on OWL DL for maximum expressiveness, but which violate the constraints of Description Logic reasoners.



OWL2 profiles (sublanguages), each profile is more restrictive than OWL DL:

- **OWL 2 EL** enables polynomial time algorithms for all the standard reasoning tasks
- OWL 2 QL enables conjunctive queries to be answered in LogSpace using standard relational database technology
- OWL 2 RL enables the implementation of polynomial time reasoning algorithms using rule-extended database technologies operating directly on RDF triples



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- Datatypes (\mathcal{D})
- ABox: class and property assertions e.g., Student(ST001), hasParent(MARIA, NIKOS), instance equality e.g., ST001 = NIKOS, and difference e.g., MARIA ≠ NIKOS

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- **ABox:** negated property assertions e.g., \neg hasPet(NIKOS, PLUTO)

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- The ontology itself also has an IRI, it can be annotated (owl:version, rdfs:label, rdfs:seeAlso, etc) and reused (owl:import)

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- properties: distinguished to ObjectProperty and DataProperty (the former relates instances to each other, the latter relates instances to data values),
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- Expressions: describing complex classes of elements in the domain (i.e. complex concepts or roles in DL terms).
- Axioms are statements that are asserted to be true (e.g., a subclass axiom)
- DL reasoners can be employed to draw inferences from asserted knowledge

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- Turtle

$Parent \equiv \exists hasChild.Person$

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Example (OWL/XML Syntax)

$Parent \equiv \exists hasChild.Person$

```
. . .
<Declaration>
    <Class IRI="#Parent"/>
</Declaration>
<Declaration>
    <Class IRI="#Person"/>
</Declaration>
<Declaration>
    <ObjectProperty IRI="#hasChild"/>
</Declaration>
<EquivalentClasses>
    <Class IRI="#Parent"/>
    <ObjectSomeValuesFrom>
        <ObjectProperty IRI="#hasChild"/>
        <Class IRI="#Person"/>
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Example (Functional/Manchester Syntax)

 $Parent \equiv \exists hasChild.Person$

Functional Syntax:

Manchester Syntax:

```
ObjectProperty: <http://example.gr#hasChild>
Class: <http://example.gr#Person>
Class: <http://example.gr#Parent>
    EquivalentTo:
        <http://example.gr#hasChild> some
        <http://example.gr#Person>
```

$Parent \equiv \exists hasChild.Person$

Turtle Syntax:

```
:Person rdf:type owl:Class .
:hasChild rdf:type owl:ObjectProperty .
:Parent rdf:type owl:Class ;
    owl:equivalentClass [ rdf:type owl:Restriction ;
        owl:onProperty :hasChild ;
        owl:someValuesFrom :Person
] .
```

OWL classes represent sets of individuals. Predefined classes in OWL are:

- owl:Thing, which represents the set of all individuals
- owl:Nothing, which represents the empty set

OWL ObjectProperties connect pairs of individuals, e.g.,

```
:parentOf(:Homer :Bart)
```

Built-in object properties:

- owl:topObjectProperty, which connects all possible pairs of individuals.
- owl:bottomObjectProperty, which does not connect any pair of individuals.

Object properties can be used to form object property expressions.

An inverse object property expression ObjectInverseOf(P) connects an individual I1 with I2 if and only if the object property P connects I2 with I1. **Example:** if an ontology contains:

```
ObjectPropertyAssertion(:parentOf :Homer :Bart)
then it also entails:
```

ObjectPropertyAssertion(ObjectInverseOf(:fatherOf)

```
:Bart :Homer)
```

In some cases, a property and its inverse coincide, or in other words, the direction of a property doesn't matter.

- Example: if an ontology contains:
- SymmetricObjectProperty(:hasSpouse)

```
ObjectPropertyAssertion(:hasSpouse :Bart :Lisa)
then it also entails:
```

ObjectPropertyAssertion(:hasSpouse :Lisa :Bart)

A property can also be asymmetric, if it connects A with B, but it never connects B with A.

Example:

AsymmetricObjectProperty(:hasChild)

Reflexive properties relate everything to itself.

Example: Everyone is a relative to him/herself:

ReflexiveObjectProperty(:hasRelative)

Note: this does not necessarily mean that every two individuals which are related by a reflexive property are identical.

Irreflexive properties model the case where no individual can be related to itself by such a property.

Example:

IrreflexiveObjectProperty(:parentOf)
Nobody can be his own parent.

Some properties relate a subject to at most one object. These properties are called functional.

- **Example:** If an ontology contains
- FunctionalObjectProperty(:hasHusband)
- ObjectPropertyAssertion(:hasHusband :Marge :Homer)

ObjectPropertyAssertion(:hasHusband :Marge :HomerSimpson)
it also entails:

- SameIndividual(:Homer :HomerSimpson)
- Note: this expression does not require every individual to have a husband, it only states that there can be no more than one.

It is also possible to indicate that the inverse of a given property is functional.

- **Example:** If an ontology contains
- InverseFunctionalObjectProperty(:hasHusband)
- ObjectPropertyAssertion(:hasHusband :Marge :Homer)

ObjectPropertyAssertion(:hasHusband :MargeBouvier :Homer)
it also entails:

ObjectPropertyAssertion(owl:sameAs :Marge :MargeBouvier) This indicates that if two or more individuals are related with the same individual via an inverse functional property, then these individuals refer to the same entity in the domain. A transitive property R interlinks an individual i with j, whenever R relates i with k, and k with j.

Example: If an ontology contains

TransitiveObjectProperty(:hasAncestor)

ObjectPropertyAssertion(:hasAncestor :Marge :ClancyBouvier)

ObjectPropertyAssertion(:hasAncestor :Lisa :Marge) it also entails:

ObjectPropertyAssertion(:hasAncestor :Lisa :ClancyBouvier)

Data properties (e.g., :hasAge) connect individuals with literals. Built-in properties:

- owl:topDataProperty, which connects all possible individuals with all literals.
- owl:bottomDataProperty, which does not connect any individual with a literal.

Annotation properties can be used to provide an annotation for an ontology, axiom, or a resource. Users can define their own annotation properties or use the available built-in annotation properties:

- rdfs:label, rdfs:comment, rdfs:seeAlso, rdfs:isDefinedBy
- owl:deprecated, owl:versionInfo, owl:priorVersion, owl:backwardCompatibleWith, owl:incompatibleWith

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- Named individuals are given an explicit name (an IRI e.g., :Peter) that can be used in any ontology to refer to the same object.
- Anonymous individuals do not have a global name. They can be defined using a name (e.g., _:somebody) local to the ontology they are contained in. They are like blank nodes in RDF.

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- rdfs:Literal contains all the elements of other data types.
- There are also the OWL datatypes owl:real and owl:rational.

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 - The value space is the set of values of the datatype. Elements of the value space are called data values.
 - The lexical space is a set of strings that can be used to refer to data values. Each member of the lexical space is called a lexical form, and it is mapped to a particular data value.

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 - The lexical space is a set of strings that can be used to refer to data values. Each member of the lexical space is called a lexical form, and it is mapped to a particular data value.
 - The facet space is a set of pairs of the form (F,v) where F is an IRI called a constraining facet, and v is an arbitrary data value called the constraining value. Each such pair is mapped to a subset of the value space of the datatype.

We can define a new datatype for a person's age by constraining the datatype integer to values between (inclusively) 0 and 120.

```
DatatypeDefinition( :personAge
DatatypeRestriction(
    xsd:integer
        xsd:minInclusive "0"^^xsd:integer
        xsd:maxInclusive "120"^^xsd:integer
    )
```

Literals represent data values such as particular strings or integers. They are analogous to RDF literals.

Examples:

"1"^^xsd:integer (typed literal) "Family Guy" (plain literal, an abbreviation for "Family Guy"^^rdf:PlainLiteral) "Padre de familia"@es (plain literal with language tag) • Class names and property expressions can be used to construct class expressions.

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- These are essentially the complex concepts or descriptions that we can define in DLs.
- Class expressions represent sets of individuals by formally specifying conditions on the individuals' properties; individuals satisfying these conditions are said to be instances of the respective class expressions.

Class expressions can be formed by:

• Applying the standard Boolean connectives to simpler class expressions or by enumerating the individuals that belong to an expression.

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```
An intersection class expression

DbjectIntersectionOf(CE_1 \dots CE_n) contains all individuals that

are instances of all class expressions CE_i for 1 \le i \le n.

Example:
```

ObjectIntersectionOf(:Dog :CanTalk)

```
A union class expression <code>ObjectUnionOf(CE_1 ... CE_n)</code> contains all individuals that are instances of at least one class expression <code>CE_i</code> for 1 \le i \le n.

Example:

<code>ObjectUnionOf(:Man :Woman)</code>
```

A complement class expression <code>ObjectComplementOf(CE)</code> contains all individuals that are not instances of the class expression CE. **Example:**

```
ObjectComplementOf( :Man )
```

We can define disjoint classes, i.e. those sets that cannot have a common element.

Example: If an ontology contains:

DisjointClasses(:Man :Woman)

ClassAssertion(:Woman :Marge)

it also entails:

ClassAssertion(ObjectComplementOf(:Man) :Marge)

An enumeration of individuals ObjectOneOf(a_1 ... a_n) contains exactly the individuals a_i with $1 \le i \le n$. Example: ObjectOneOf(:Saturday :Sunday)

From

EquivalentClasses(:GriffinFamilyMember

ObjectOneOf(:Peter :Lois :Stewie :Meg :Chris :Brian)) DifferentIndividuals(:Quagmire :Peter :Lois :Stewie :Meg

:Chris :Brian)

we can infer ClassAssertion(

ObjectComplementOf(:GriffinFamilyMember) :Quagmire)

From

```
ClassAssertion(:GriffinFamilyMember :Peter)
ClassAssertion(:GriffinFamilyMember :Lois)
ClassAssertion(:GriffinFamilyMember :Stewie)
ClassAssertion(:GriffinFamilyMember :Meg)
ClassAssertion(:GriffinFamilyMember :Chris)
ClassAssertion(:GriffinFamilyMember :Brian)
```

```
DifferentIndividuals(:Quagmire :Peter :Lois :Stewie
:Meg :Chris :Brian)
```

Can we infer this: ClassAssertion(ObjectComplementOf(:GriffinFamilyMember) :Quagmire)? An existential class expression ObjectSomeValuesFrom(OPE CE) consists of an object property expression OPE and a class expression CE, and it contains all those individuals that are connected by OPE to an individual that is an instance of CE.

Example:

ObjectSomeValuesFrom(:hasChild :Person)

If OPE is simple, the above class expression is equivalent with the class expression ObjectMinCardinality(1 OPE CE)

From

```
ObjectPropertyAssertion(:hasChild :Peter :Stewie)
ClassAssertion(:Person :Stewie)
```

we can infer

```
ClassAssertion(
ObjectSomeValuesFrom(:hasChild :Person) :Peter)
```

A universal class expression ObjectAllValuesFrom(OPE CE) consists of an object property expression OPE and a class expression CE, and it contains all those individuals that are connected by OPE only to individuals that are instances of CE.

Example:

ObjectAllValuesFrom(:fatherOf :Man)

If OPE is simple, the above class expression is equivalent with the class expression ObjectMaxCardinality(0 OPE ObjectComplementOf(CE)) An individual value class expression ObjectHasValue(OPE a) consists of an object property expression OPE and an individual a, and it contains all those individuals that are connected by OPE to a.

Example: If an ontology contains:

EquivalentClasses(:SolarPlanet ObjectHasValue(:orbits :Sun)) ObjectPropertyAssertion(:orbits :Earth :Sun))

it can also entail:

ClassAssertion(:SolarPlanet :Earth)

The above class expression is equivalent to the class expression ObjectSomeValuesFrom(OPE ObjectOneOf(a)).

A self-restriction ObjectHasSelf(OPE) consists of an object property expression OPE, and it contains all those individuals that are connected by OPE to themselves.

Example: if an ontology contains:

EquivalentClasses(:Narcisist ObjectHasSelf(:likes))

ObjectPropertyAssertion(:likes :Peter :Peter)
it also infers:

ClassAssertion(:Narcisist :Peter)

Object property cardinality restrictions are distinguished into:

- Qualified: apply only to individuals that are connected by the object property expression and are instances of the qualifying class expression. (e.g. >3hasChild.Male)
- Unqualified: apply to all individuals that are connected by the object property expression (this is equivalent to the qualified case with the qualifying class expression equal to owl:Thing) (e.g. >3hasChild).

A minimum cardinality expression ObjectMinCardinality(n OPE CE) consists of a nonnegative integer n, an object property expression OPE, and a class expression CE, and it contains all those individuals that are connected by OPE to at least n different individuals that are instances of CE. If CE is missing, it is taken to be owl:Thing.

Example:

ObjectMinCardinality(2 :fatherOf :Man)

From

```
ObjectPropertyAssertion(:fatherOf :Peter :Stewie)
ObjectPropertyAssertion(:fatherOf :Peter :Chris)
ClassAssertion(:Man :Stewie)
ClassAssertion(:Man :Chris)
DifferentIndividuals(:Chris :Stewie)
```

we can infer: ClassAssertion(ObjectMinCardinality(2 :fatherOf :Man) :Peter) A maximum cardinality expression ObjectMaxCardinality(n OPE CE) consists of a nonnegative integer n, an object property expression OPE, and a class expression CE, and it contains all those individuals that are connected by OPE to at most n different individuals that are instances of CE. If CE is missing, it is taken to be owl:Thing.

Example:

ObjectMaxCardinality(2 :hasPet)

An exact cardinality expression ObjectExactCardinality(n OPE CE) consists of a nonnegative integer n, an object property expression OPE, and a class expression CE, and it contains all those individuals that are connected by OPE to exactly n different individuals that are instances of CE. **Example:**

ObjectExactCardinality(1 :hasPet :Dog)

The above expression is equivalent to

ObjectIntersectionOf(

ObjectMinCardinality(n OPE CE)

ObjectMaxCardinality(n OPE CE))

- Data property restrictions are similar to the restrictions on object property expressions.
- The main difference is that the expressions for existential and universal quantification allow for n-ary data ranges.
- Given the syntax for data ranges given earlier, only unary data ranges are supported.
- However, the specification provide the syntactic constructs needed to have n-ary data ranges e.g., sets of rectangles defined by appropriate geometric constraints.

The "Data Range Extension: Linear Equations" W3C note proposes an extension to OWL 2 for defining n-ary data ranges in terms of linear (in)equations with rational coefficients. See http://www.w3.org/TR/owl2-dr-linear/.

A subclass axiom SubClassOf(CE1 CE2) states that the class expression CE1 is a subclass of the class expression CE2.

Example:

SubClassOf(:Child :Person)

The properties known from RDFS for SubClassOf hold here as well (Reflexivity, Transitivity)

An equivalent classes axiom EquivalentClasses(CE_1 ... CE_n) states that all of the class expressions CE_i, $1 \le i \le n$, are semantically equivalent to each other.

Example:

EquivalentClasses(:Boy ObjectIntersectionOf(:Child :Male)) An axiom EquivalentClasses(CE1 CE2) is equivalent to the conjunction of the following two axioms: SubClassOf(CE1 CE2) SubClassOf(CE2 CE1) A disjoint classes axiom DisjointClasses(CE_1 ... CE_n) states that all of the class expressions CE_i, $1 \le i \le n$, are pairwise disjoint. Example:

```
DisjointClasses(:Boy :Girl)
```

An axiom DisjointClasses(CE1 CE2) is equivalent to the following axiom:

SubClassOf(CE1 ObjectComplementOf(CE2))

A disjoint union axiom DisjointUnion(C CE_1 ... CE_n) states that a class C is a disjoint union of the class expressions CE_i, $1 \le i \le n$, all of which are pairwise disjoint.

Such axioms are sometimes referred to as covering axioms, as they state that the extensions of all CE_i exactly cover the extension of C. Example:

DisjointUnion(:Child :Boy :Girl)

Each such axiom is equivalent to the conjunction of the following two axioms: EquivalentClasses(C ObjectUnionOf(CE1 ... CEn)) DisjointClasses(CE1 ... CEn)

From

```
DisjointUnion(:Child :Boy :Girl)
ClassAssertion(:Child :Stewie)
ClassAssertion(ObjectComplementOf(:Girl) :Stewie)
```

we can infer ClassAssertion(:Boy :Stewie)

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- Object subproperty axioms are analogous to subclass axioms.
- The basic form of an object subproperty axiom is SubObjectPropertyOf(OPE1 OPE2).
- This axiom states that the object property expression OPE1 is a subproperty of the object property expression OPE2 i.e. if an individual x is connected by OPE1 to an individual y, then x is also connected by OPE2 to y.
- SubObjectPropertyOf is a reflexive and transitive relation.

If OPE1, ..., OPEn are object properties then
OPE1 ... OPEn is called an object property chain.
The more complex form of object subproperty axioms is
SubObjectPropertyOf(
ObjectPropertyChain(OPE1 ... OPEn) OPE).
This axiom states that, if an individual x1 is connected by a sequence of
object property expressions OPE1, ..., OPEn with an individual xn, then
x1 is also connected with xn by the object property expression OPE.
These axioms are known as complex role inclusions in the DL literature.

From

SubObjectPropertyOf(
ObjectPropertyChain(:hasMother :hasSister) :hasAunt)

ObjectPropertyAssertion(:hasMother :Stewie :Lois)
ObjectPropertyAssertion(:hasSister :Lois :Carol)

we can infer
ObjectPropertyAssertion(:hasAunt :Stewie :Carol)

An equivalent object properties axiom

EquivalentObjectProperties(OPE_1 ... OPE_n) states that all of the object property expressions OPE_i, $1 \le i \le n$, are semantically equivalent to each other.

The axiom EquivalentObjectProperties(OPE1 OPE2) is equivalent to the following two axioms:

SubObjectPropertyOf(OPE1 OPE2)

SubObjectPropertyOf(OPE2 OPE1)

A disjoint object properties axiom DisjointObjectProperties(OPE1 ... OPEn) states that all of the object property expressions OPE_i, $1 \le i \le n$, are pairwise disjoint. **Example:**

DisjointObjectProperties(:hasFather :hasMother)

- An object property domain axiom ObjectPropertyDomain(OPE CE) states that the domain of the object property expression OPE is the class expression CE i.e. if an individual x is connected by OPE with some other individual, then x is an instance of CE.
- An object property range axiom ObjectPropertyRange(OPE CE) states that the range of the object property expression OPE is the class expression CE i.e. if some individual is connected by OPE with an individual x, then x is an instance of CE.

• Individual equality

- Individual equality
- Individual inequality

- Individual equality
- Individual inequality
- Class assertion

- Individual equality
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- Class assertion
- Positive object property assertion

- Individual equality
- Individual inequality
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- Positive object property assertion
- Negative object property assertion

- Individual equality
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- Negative object property assertion
- Positive data property assertion

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- Negative data property assertion

Assertions are often also called facts. They are part of the ABox in DLs.

An individual equality axiom SameIndividual(a_1 ... a_n) states that all of the individuals a_i, $1 \le i \le n$, are equal to each other. Example:

From SameIndividual(:Meg :Megan)
ObjectPropertyAssertion(:hasBrother :Meg :Stewie)
we can infer:

ObjectPropertyAssertion(:hasBrother :Megan :Stewie)

An individual inequality axiom DifferentIndividuals(a_1 ... a_n) states that all of the individuals a_i, $1 \le i \le n$, are different from each other.

Example:

DifferentIndividuals(:Peter :Meg :Chris :Stewie)

A class assertion ClassAssertion(CE a) states that the individual a is an instance of the class expression CE.

Example:

ClassAssertion(:Dog :Brian)

A positive object property assertion

ObjectPropertyAssertion(OPE a_1 a_2) states that the individual a_1 is connected by the object property expression OPE to the individual a_2.

A negative object property assertion

NegativeObjectPropertyAssertion(OPE a_1 a_2) states that the individual a_1 is not connected by the object property expression OPE to the individual a_2.

Examples:

ObjectPropertyAssertion(:hasDog :Peter :Brian)
NegativeObjectPropertyAssertion(:hasSon :Peter :Meg)

Readings

- OWL 2 Web Ontology Language Primer (Second Edition) https://www.w3.org/TR/owl2-primer/
- OWL 2 Web Ontology Language Manchester Syntax (Second Edition) https://www.w3.org/TR/owl2-manchester-syntax/
- Krötzsch, M. (2012). OWL 2 Profiles: An Introduction to Lightweight Ontology Languages. In: Eiter, T., Krennwallner, T. (eds) Reasoning Web. Semantic Technologies for Advanced Query Answering. Reasoning Web 2012. Lecture Notes in Computer Science, vol 7487. Springer, Berlin, Heidelberg. https://doi.org/10.1007/978-3-642-33158-9_4
- An introduction to OWL 2 and DL SROIQ, which also discusses various OWL tools/applications: Ian Horrocks and Peter F. Patel-Schneider. KR and Reasoning on the Semantic Web: OWL. In Handbook of Semantic Web Technologies, chapter 9. Springer, 2010. http://www.cs.ox.ac.uk/people/ian.horrocks/Publications/download/2010/HoPa10a.pdf